

# Reverse triple I method of restriction for fuzzy reasoning\*

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Received August 14, 2001; revised October 8, 2001

**Abstract** A theory of reverse triple I method of restriction for implication operator  $R_0$  is proposed. And the general computation formulas of infimum for fuzzy modus ponens and supremum for fuzzy modus tollens of  $\alpha$ -reverse triple I method of restriction are obtained respectively.

**Keywords:** fuzzy reasoning, implication operator  $R_0$ , reverse triple I method of restriction.

Since Zadeh<sup>[1]</sup> proposed the method of compositional rule of inference (CRI) for fuzzy reasoning, the method has been generalized respectively from diversified aspects in Refs. [2 ~ 4]. Generally, because this method does not have the reversibility properties, the MP-approximation property of the method of CRI was studied using certain common compositional operators and implication operators in Ref. [5]. Moreover, in order to avoid some drawbacks of the method of CRI, Wang<sup>[6]</sup> proposed first triple I method with total inference rule that utilizes the implication operator in every step of the reasoning. Afterwards, the theory of restriction degree of triple I method was further presented by the authors<sup>[7]</sup>, its generalization form should be expressed as the following optimal problem.

For any  $\alpha \in (0, 1]$ ,  $A \in \mathcal{F}(X)$ ,  $B \in \mathcal{F}(Y)$ , and  $A^* \in \mathcal{F}(X)$  (or  $B^* \in \mathcal{F}(Y)$ ), seek the optimal  $B^* \in \mathcal{F}(Y)$  (or  $A^* \in \mathcal{F}(X)$ ) satisfying

$$(A(x) \rightarrow B(y)) \rightarrow (A^*(x) \rightarrow B^*(y)) \leq \alpha, \tag{1}$$

for any  $x \in X$  and  $y \in Y$ , where  $\mathcal{F}(X)$  and  $\mathcal{F}(Y)$  denote respectively the collections consisting of all fuzzy subsets of  $X$  and  $Y$ . And for the implication operator  $R_0: [0, 1]^2 \rightarrow [0, 1]$ :

$$R_0(a, b) = \begin{cases} 1, & a \leq b, \\ a' \vee b, & a > b, \end{cases} \quad (\text{here } a' = 1 - a).$$

General computation formulas of supremum for fuzzy modus ponens (FMP) and infimum for fuzzy modus tollens (FMT) of  $\alpha$ -triple I method are obtained respectively in Ref. [7].

In this paper, the theory of reverse triple I method of restriction is proposed, its generalization form should be represented as the follows.

Under the hypotheses of (1), seek the optimal  $B^* \in \mathcal{F}(Y)$  (or  $A^* \in \mathcal{F}(X)$ ) such that

$$(A^*(x) \rightarrow B^*(y)) \rightarrow (A(x) \rightarrow B(y)) \leq \alpha, \tag{2}$$

for any  $x \in X$  and  $y \in Y$ .

The computation formulas of infimum for FMP and supremum for FMT of  $\alpha$ -reverse triple I method of restriction are given respectively, using the implication operator  $R_0$  in this paper.

## 1 Infimum for FMP of $\alpha$ -reverse triple I method of restriction

Now, we consider the generalization problems of reverse triple I method of restriction, i.e. for given  $\alpha \in (0, 1]$ , seek the optimal solution satisfying (2). At first, for the generalization problem of FMP, we give the following principle of restriction of  $\alpha$ -reverse triple I method.

**Principle of restriction for  $\alpha$ -reverse triple I FMP.** Suppose that  $X$  and  $Y$  are non-empty sets,  $A, A^* \in \mathcal{F}(X)$ ,  $B \in \mathcal{F}(Y)$ . If  $B^*$  is the minimal fuzzy set in  $\mathcal{F}(Y)$  to satisfy (2), then  $B^*$  is called the  $\alpha$ -solution of (2) for the reverse triple I FMP.

**Remark 1.** For any  $y \in Y$  when  $B^*(y) \equiv 1$ , the left side of (2) will always take its minimal value  $R_0(A(x), B(y))$ ; when  $B^*(y) \equiv 0$ , it will always

\* Supported by the National Natural Science Foundation of China (Grant Nos. 60074015, 60004010) and Research Fund at Tsinghua University (Grant No. JC2001029)

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take its maximal value

$$(A^*(x))' \rightarrow R_0(A(x), B(y))$$

$$= \begin{cases} 1, & (A^*(x))' \leq R_0(A(x), B(y)), \\ A^*(x), & (A^*(x))' > R_0(A(x), B(y)), A^*(x) > R_0(A(x), B(y)), \\ R_0(A(x), B(y)), & (A^*(x))' > R_0(A(x), B(y)), A^*(x) \leq R_0(A(x), B(y)). \end{cases} \quad (3)$$

By (3), when

$$(A^*(x))' > R_0(A(x), B(y))$$

and

$$A^*(x) \leq R_0(A(x), B(y)), \quad (4)$$

the maximal value of the left side of (2) is  $R_0(A(x), B(y))$ , and its minimal value is always  $R_0(A(x), B(y))$ . So, if (4) holds, for any  $\alpha \in [R_0(A(x), B(y)), 1]$ , the minimal fuzzy set in  $\mathcal{F}(Y)$  satisfying (2) is  $B^*(y) \equiv 0$ . If (4) does not hold, for the generalization problem of FMP, the range of  $\alpha$  should be confined as

$$\alpha \in (R_0(A(x), B(y)), 1). \quad (5)$$

Furthermore, we have the formula of infimum for FMP of  $\alpha$ -reverse triple I method of restriction as follows.

**Theorem 1. (formula of infimum for FMP of  $\alpha$ -reverse triple I method of restriction 1)** Suppose that  $X, Y$  are non-empty sets,  $A$  and  $A^* \in \mathcal{F}(X)$  and  $B \in \mathcal{F}(Y)$ . For any  $y \in Y$ , if  $x \in E_y$ , then the infimum  $B^*(y)$  consisting of fuzzy sets in  $\mathcal{F}(Y)$  to satisfy (2) is determined by

$$B^*(y) = \sup_{x \in E_y \cap K_y} [A^*(x) \wedge (R_0(A(x), B(y)) \vee \alpha')] \cdot \chi_{E_y \cap K_y} + \sup_{x \in E_y - K_y} [A^*(x) \wedge (R_0(A(x), B(y)))] \cdot \chi_{E_y - K_y}, \quad (6)$$

where  $K_y = \{x \in X \mid A^*(x) > \alpha\}$  and  $E_y = \{x \in X \mid (A^*(x))' \leq R_0(A(x), B(y))\}$ .

**Proof.** For any  $y \in Y$ , we will first prove: any  $C(y) \in \mathcal{F}(Y)$  with  $C(y) > B^*(y)$  must satisfy (2). For  $x \in E_y$ , we will discuss in different cases as follows.

**Case 1.** If  $x \in E_y \cap K_y$ , then using  $B^*(y)$  determined by (6), we have

$$C(y) > B^*(y) \geq A^*(x) \wedge (R_0(A(x), B(y)) \vee \alpha'). \quad (7)$$

Further, the discussion will be partitioned again into two possible cases.

(i) If  $A^*(x) \leq C(y)$  then  $R_0(A^*(x), C(y)) \equiv 1$ , consequently

$$M_{xy} = R_0(A^*(x), C(y)) \rightarrow R_0(A(x), B(y)) = 1 \rightarrow R_0(A(x), B(y)) = R_0(A(x), B(y)) \leq \alpha.$$

(ii) If  $A^*(x) > C(y)$ , then  $R_0(A^*(x), C(y)) = (A^*(x))' \vee C(y)$ . By (7), we have  $C(y) > B^*(y) \geq R_0(A(x), B(y)) \vee \alpha'$ . (8)

This implies that  $(C(y))' < \alpha$ . Moreover, by the hypotheses in (8) and (5), we have

$$M_{xy} = R_0(A^*(x), C(y)) \rightarrow R_0(A(x), B(y)) = (A^*(x) \wedge (C(y))') \vee R_0(A(x), B(y)) < \alpha. \quad (9)$$

**Case 2.** If  $x \in E_y - K_y$ , similarly, from the sense of  $B^*(y)$  determined by (6), we know  $C(y) > B^*(y) \geq A^*(x) \wedge R_0(A(x), B(y))$ . (10)

The discussion will be partitioned into two possible cases.

(i) If  $A^*(x) \leq C(y)$ , then from the proof of (i) in Case 1 of Theorem 1, we know that  $C(y)$  satisfies (2).

(ii) If  $A^*(x) > C(y)$ , then  $R_0(A^*(x), C(y)) = (A^*(x))' \vee C(y)$ . From (10), we have  $C(y) > B^*(y) \geq R_0(A(x), B(y))$ . (11)

From this, and noting that  $x \notin K_y$ , we deduce  $A^*(x) \leq \alpha$ . In addition, using the hypothesis in (5), we get

$$M_{xy} = R_0(A^*(x), C(y)) \rightarrow R_0(A(x), B(y)) = (A^*(x) \wedge (C(y))') \vee R_0(A(x), B(y)) \leq \alpha. \quad (12)$$

Combining the proofs of Case 1 and Case 2, it follows that  $C(y)$  satisfies (2). On the other hand, we will prove: for some  $y_0 \in Y$  and any  $D(y_0) \in \mathcal{F}(Y)$  with  $D(y_0) < B^*(y_0)$ ,  $D(y_0)$  cannot satisfy (2). In fact, applying the sense of  $B^*(y_0)$  determined by (6), the discussion will be partitioned into two possible cases.

**Case 1.** There exists  $x_0 \in E_{y_0} \cap K_{y_0}$  such that  $D(y_0) < A^*(x_0) \wedge (R_0(A(x_0), B(y_0)) \vee \alpha')$ . (13)

We will discuss again in two possible cases.

(i) If  $D(y_0) \leq R_0(A(x_0), B(y_0))$ . Then  $R_0(A^*(x_0), D(y_0)) \rightarrow R_0(A(x_0), B(y_0)) = ((A^*(x_0))' \vee D(y_0)) \rightarrow R_0(A(x_0), B(y_0)) = 1 > \alpha$ .

(ii) If  $D(y_0) > R_0(A(x_0), B(y_0))$ , then from (13), it is known that  $D(y_0) < \alpha'$ . Noting that  $x_0 \in K_{y_0}$ , we have

$$\begin{aligned} R_0(A^*(x_0), D(y_0)) &\rightarrow R_0(A(x_0), B(y_0)) \\ &= ((A^*(x_0))' \vee D(y_0)) \rightarrow R_0(A(x_0), B(y_0)) \\ &= (A^*(x_0) \wedge (D(y_0))' \vee R_0(A(x_0), B(y_0))) \\ &> \alpha. \end{aligned}$$

**Case 2.** There exists  $x_0 \in E_{y_0} - K_{y_0}$  such that  $D(y_0) < A^*(x_0) \wedge R_0(A(x_0), B(y_0))$ . (14)

Consequently, we have

$$\begin{aligned} R_0(A^*(x_0), D(y_0)) &\rightarrow R_0(A(x_0), B(y_0)) \\ &= ((A^*(x_0))' \vee D(y_0)) \rightarrow R_0(A(x_0), B(y_0)) \\ &= 1 > \alpha. \end{aligned} \quad (15)$$

So,  $D(y_0)$  cannot satisfy (2). All of these show that  $B^*(y)$  is the infimum consisting of fuzzy sets in  $\mathcal{F}(Y)$  to satisfy (2).

**Theorem 2. (formula of infimum for FMP of  $\alpha$ -reverse triple I method of restriction 2)** Suppose that  $X, Y$  are non-empty sets,  $A$  and  $A^* \in \mathcal{F}(X)$  and  $B \in \mathcal{F}(Y)$ . For any  $y \in Y$ , if  $x \in F_y$ , then the infimum  $B^*(y)$  consisting of fuzzy sets in  $\mathcal{F}(Y)$  to satisfy (2) is determined by

$$B^*(y) = \sup_{x \in F_y \cap K_y} [A^*(x) \wedge \alpha']. \quad (16)$$

Here  $K_y = \{x \in X \mid A^*(x) > \alpha\}$  and  $F_y = \{x \in X \mid A^*(x) \wedge (A^*(x))' > R_0(A(x), B(y))\}$ .

**Proof.** For any  $y \in Y$  and  $B^*(y)$  determined by (16), we will first prove: any  $C(y) \in \mathcal{F}(Y)$  with  $C(y) > B^*(y)$  must satisfy (2). For  $x \in F_y$ , we will discuss in different cases as follows.

**Case 1.** If  $x \in F_y \cap K_y$ , then applying  $B^*(y)$  determined by (16), we know

$$C(y) > B^*(y) \geq A^*(x) \wedge \alpha'. \quad (17)$$

We will discuss again in two possible cases.

(i) If  $A^*(x) \leq C(y)$ , then from the proof of

(i) in Case 1 of Theorem 1, we know that  $C(y)$  satisfies (2).

(ii) If  $A^*(x) > C(y)$ , then  $R_0(A^*(x), C(y)) = (A^*(x))' \vee C(y)$ . Noting (17), we have  $C(y) > B^*(y) \geq \alpha'$ . It follows that  $(C(y))' < \alpha$ . By the hypothesis in Case 1 of Theorem 2, and the hypothesis in (5), we have

$$\begin{aligned} M_{xy} &= R_0(A^*(x), C(y)) \rightarrow R_0(A(x), B(y)) \\ &= (A^*(x) \wedge (C(y))') \vee R_0(A(x), B(y)) \\ &< \alpha. \end{aligned} \quad (18)$$

**Case 2.** If  $x \in F_y - K_y$ , similarly, from the sense of  $B^*(y)$  determined by (16), we know  $B^*(y) = 0$ . Noting that  $x \in F_y$ ,  $x \notin K_y$ , and the hypothesis in (5), we have

$$\begin{aligned} M_{xy} &= R_0(A^*(x), B^*(y)) \rightarrow R_0(A(x), B(y)) \\ &= (A^*(x))' \rightarrow R_0(A(x), B(y)) \\ &= A^*(x) \vee R_0(A(x), B(y)) \leq \alpha. \end{aligned} \quad (19)$$

Combining the proofs of Case 1 and Case 2, we know that  $C(y)$  satisfies (2). On the other hand, for some  $y_0 \in Y$  with  $B^*(y_0) > 0$ , any  $D(y_0)$  provided  $D(y_0) < B^*(y_0)$  cannot satisfy (2). In fact, by the sense of  $B^*(y_0)$  determined by (16), there exists  $x_0 \in F_{y_0} \cap K_{y_0}$  such that  $D(y_0) < A^*(x_0) \wedge \alpha'$ . So, noting that  $x_0 \in F_{y_0}$  and  $x_0 \in K_{y_0}$ , and the hypothesis in (5), we obtain

$$\begin{aligned} R_0(A^*(x_0), D(y_0)) &\rightarrow R_0(A(x_0), B(y_0)) \\ &= ((A^*(x_0))' \vee D(y_0)) \rightarrow R_0(A(x_0), B(y_0)) \\ &= (A^*(x_0) \wedge (D(y_0))' \vee R_0(A(x_0), B(y_0))) \\ &> \alpha. \end{aligned}$$

So,  $D(y_0)$  cannot satisfy (2). Sum up the above proof,  $B^*(y) \in \mathcal{F}(Y)$  is the infimum consisting of fuzzy sets to satisfy (2).

## 2 Supremum for FMT of $\alpha$ -reverse triple I method of restriction

Now, let us consider the generalization problem of FMT, we give the following principle of restriction of the  $\alpha$ -reverse triple I method.

**Principle of restriction for  $\alpha$ -reverse triple I FMT.** Suppose that  $X$  and  $Y$  are non-empty sets,  $A \in \mathcal{F}(X)$ ,  $B$  and  $B^* \in \mathcal{F}(Y)$ . If  $A^*$  is the maximal fuzzy set in  $\mathcal{F}(X)$  to satisfy (2), then  $A^*$  is called the  $\alpha$ -solution of (2) for the reverse triple I FMT.

**Remark 2.** For any  $x \in X$ , when  $A^*(x) \equiv 0$ , the left side of (2) will always take its minimal value

$R_0(A(x), B(y))$ ; and when  $A^*(x) \equiv 1$ , it will always take its maximal value

$$B^*(y) \rightarrow R_0(A(x), B(y)) = \begin{cases} 1, & B^*(y) \leq R_0(A(x), B(y)), \\ R_0(A(x), B(y)), & B^*(y) > R_0(A(x), B(y)), (B^*(y))' \leq R_0(A(x), B(y)), \\ (B^*(y))', & B^*(y) > R_0(A(x), B(y)), (B^*(y))' > R_0(A(x), B(y)). \end{cases} \quad (20)$$

By (20), when

$$(B^*(y))' > R_0(A(x), B(y))$$

and

$$(A^*(y))' \leq R_0(A(x), B(y)), \quad (21)$$

the maximal value of the left side of (2) is  $R_0(A(x), B(y))$ , and its minimal value is always  $R_0(A(x), B(y))$ . Hence, if (21) holds, for any  $\alpha \in [R_0(A(x), B(y)), 1]$ , the maximal fuzzy set in  $\mathcal{F}(X)$  satisfying (2) is  $A^*(x) \equiv 1$ .

If (21) does not hold, for the generalization problem of FMT, the range of  $\alpha$  should also be confined by (5).

Furthermore, we have the formula of supremum for FMT of  $\alpha$ -reverse triple I method of restriction as follows.

**Theorem 3. (formula of supremum for FMT of  $\alpha$ -reverse triple I method of restriction 1)** Suppose that  $X, Y$  are non-empty sets,  $A \in \mathcal{F}(X)$ ,  $B$  and  $B^* \in \mathcal{F}(Y)$ . For any  $x \in X$ , if  $y \in E_x$ , then the supremum  $A^*(x)$  consisting of fuzzy sets in  $\mathcal{F}(X)$  to satisfy (2) is determined by

$$A^*(x) = \inf_{y \in E_x \cap K_x} [B^*(y) \vee (R_0'(A(x), B(y)) \wedge \alpha)] \cdot \chi_{E_x \cap K_x} + \inf_{y \in E_x - K_x} [B^*(y) \vee R_0'(A(x), B(y))] \cdot \chi_{E_x - K_x}, \quad (22)$$

where  $E_x = \{y \in Y \mid B^*(y) \leq R_0(A(x), B(y))\}$  and  $K_x = \{y \in Y \mid B^*(y) < \alpha'\}$ .

**Proof.** For any  $x \in X$  and  $C(x) < A^*(x)$ , we will first prove:  $C(x)$  must satisfy (2). For  $y \in E_x$ , we will discuss in different cases as follows.

**Case 1.** If  $y \in E_x \cap K_x$ , then from  $A^*(x)$  determined by (22), we have

$$C(x) < A^*(x) \leq B^*(y) \vee (R_0'(A(x), B(y)) \wedge \alpha). \quad (23)$$

Further, the discussion will be partitioned again into two possible cases.

(i) If  $C(x) \leq B^*(y)$ , then  $R_0(C(x), B^*(y)) = 1$ , similar to the proof of (1) in Case 1 of Theorem 1, we get

$$M_{xy} = R_0(C(x), B^*(y)) \rightarrow R_0(A(x), B(y)) \leq \alpha.$$

(ii) If  $C(x) > B^*(y)$ , then  $R_0(C(x), B^*(y)) = C'(x) \vee B^*(y)$ . Noting (23), we have  $C(x) < A^*(x) \leq R_0'(A(x), B(y)) \wedge \alpha$ .

$$(24)$$

This implies  $C(x) < \alpha$  and  $C'(x) > R_0(A(x), B(y))$ . Consequently, using the hypothesis in (5), we deduce

$$M_{xy} = R_0(C(x), B^*(y)) \rightarrow R_0(A(x), B(y)) = (C(x) \wedge (B^*(y))') \vee R_0(A(x), B(y)) < \alpha. \quad (25)$$

**Case 2.** If  $y \in E_x - K_x$ , similarly, from the sense of  $A^*(x)$  determined by (22), we have

$$C(x) < A^*(x) \leq B^*(y) \vee R_0'(A(x), B(y)). \quad (26)$$

We will discuss again in two possible cases.

(i) If  $C(x) \leq B^*(y)$ , then from (1) in Case 1 of Theorem 3,  $C(x)$  satisfies (2).

(ii) If  $C(x) > B^*(y)$ , then  $R_0(C(x), B^*(y)) = C'(x) \vee B^*(y)$ . By (26), we have  $C(x) < R_0'(A(x), B(y))$ . It follows that  $C'(x) > R_0(A(x), B(y))$ . Noting  $y \notin K_x$  and the hypothesis in (5), we have

$$M_{xy} = R_0(C(x), B^*(y)) \rightarrow R_0(A(x), B(y)) = (C(x) \wedge (B^*(y))') \vee R_0(A(x), B(y)) \leq \alpha. \quad (27)$$

Combining the discussions of Case 1 and Case 2, we know that  $C(x)$  satisfies (2). On the other hand, for some  $x_0 \in X$  and any  $D(x_0) \in \mathcal{F}(X)$  with  $D(x_0) > A^*(x_0)$ ,  $D(x_0)$  cannot satisfy (2). In fact, using the sense of  $A^*(x_0)$  determined by (22), the discussion will be partitioned again into two possible cases.

**Case 1.** There exists  $y_0 \in E_{x_0} \cap K_{x_0}$  such that

$$D(x_0) > B^*(y_0) \vee (R_0'(A(x_0), B(y_0)) \wedge \alpha). \quad (28)$$

Then, we will discuss in two possible cases.

(i) If  $D(x_0) \geq R'_0(A(x_0), B(y_0))$ , then

$$\begin{aligned} R_0(D(x_0), B^*(y_0)) &\rightarrow R_0(A(x_0), B(y_0)) \\ &= (D(x_0))' \vee B^*(y_0) \rightarrow R_0(A(x_0), B(y_0)) \\ &= 1 > \alpha. \end{aligned}$$

(ii) If  $D(x_0) < R'_0(A(x_0), B(y_0))$ , then (28) yields  $D(x_0) > \alpha$ . By  $y_0 \in K_{x_0}$ , we have

$$\begin{aligned} R_0(D(x_0), B^*(y_0)) &\rightarrow R_0(A(x_0), B(y_0)) \\ &= (D(x_0) \wedge (B^*(y_0))') \vee R_0(A(x_0), B(y_0)) \\ &> \alpha. \end{aligned}$$

**Case 2.** There exists  $y_0 \in E_{x_0} - K_{x_0}$  such that  $D(x_0) > B^*(y_0) \vee R'_0(A(x_0), B(y_0))$ . Then  $D(x_0) > R'_0(A(x_0), B(y_0))$ . By noting  $y_0 \in E_{x_0}$ , we get

$$\begin{aligned} R_0(D(x_0), B^*(y_0)) &\rightarrow R_0(A(x_0), B(y_0)) \\ &= (D(x_0))' \vee B^*(y_0) \rightarrow R_0(A(x_0), B(y_0)) \\ &= 1 > \alpha. \end{aligned}$$

So,  $D(x_0)$  cannot satisfy (2). To sum up,  $A^*(x)$  is the supremum consisting of fuzzy sets in  $\mathcal{F}(X)$  to satisfy (2).

**Theorem 4. (formula of supremum for FMT of  $\alpha$ -reverse triple I method of restriction 2)** Suppose that  $X, Y$  are non-empty sets,  $A \in \mathcal{F}(X)$ ,  $B$  and  $B^* \in \mathcal{F}(Y)$ . For any  $x \in X$ , if  $y \in F_x$ , then the supremum  $A^*(x)$  consisting of fuzzy sets in  $\mathcal{F}(X)$  to satisfy (2) is determined by

$$A^*(x) = \inf_{y \in F_x \cap K_x} [B^*(y) \vee \alpha], \quad (29)$$

where  $F_x = \{y \in Y \mid B^*(y) \wedge (B^*(y))' > R_0(A(x), B(y))\}$  and  $K_x = \{y \in Y \mid B^*(y) < \alpha'\}$ .

**Proof.** For any  $x \in X$  and  $C(x) < A^*(x)$ , we will first prove:  $C(x)$  must satisfy (2). For  $y \in F_x$  we will discuss in different cases as follows.

**Case 1.** If  $y \in F_x \cap K_x$ , then from  $A^*(x)$  determined by (29), we have

$$C(x) < A^*(x) \leq B^*(y) \vee \alpha. \quad (30)$$

We will discuss again in two possible cases.

(i) If  $C(x) \leq B^*(y)$ , then from the proof of (1) in Case 1 of Theorem 3, we know that  $C(y)$  satisfies (2).

(ii) If  $C(x) > B^*(y)$ , then  $R_0(C(x), B^*(y)) = C'(x) \vee B^*(y)$ . Noting (30), we have  $C(x) < A^*(x) \leq \alpha$ . From this, and applying the hypothesis of Case 1 of Theorem 4, and the hypothesis in (5), we get

$$\begin{aligned} M_{xy} &= R_0(C(x), B^*(y)) \rightarrow R_0(A(x), B(y)) \\ &= (C(x) \wedge (B^*(y))') \vee R_0(A(x), B(y)) \\ &< \alpha. \end{aligned} \quad (31)$$

**Case 2.** If  $y \in F_x - K_x$ , then using  $A^*(x)$  determined by (29), we have  $A^*(x) = 1$ . Noting  $y \in F_x$  and  $y \notin K_x$  the hypothesis in (5), we obtain

$$\begin{aligned} M_{xy} &= R_0(A^*(x), B^*(y)) \rightarrow R_0(A(x), B(y)) \\ &= B^*(y) \rightarrow R_0(A(x), B(y)) \\ &= (B^*(y))' \vee R_0(A(x), B(y)) \\ &= (B^*(y))' \leq \alpha. \end{aligned} \quad (32)$$

Combining the discussions of Case 1 and Case 2, we know that  $C(x)$  satisfies (2). On the other hand, for some  $x_0 \in X$  and any  $D(x_0) \in \mathcal{F}(X)$  with  $D(x_0) > A^*(x_0)$ ,  $D(x_0)$  cannot satisfy (2). In fact, using the sense of  $A^*(x_0)$  determined by (29), there exists  $y_0 \in F_{x_0} \cap K_{x_0}$  such that  $D(x_0) > B^*(y_0) \vee \alpha$ . Moreover, by  $y_0 \in F_{x_0}$  and  $y_0 \in K_{x_0}$ , we have

$$\begin{aligned} R_0(D(x_0), B^*(y_0)) &\rightarrow R_0(A(x_0), B(y_0)) \\ &= (D(x_0))' \vee B^*(y_0) \rightarrow R_0(A(x_0), B(y_0)) \\ &= (D(x_0) \wedge (B^*(y_0))') \vee R_0(A(x_0), B(y_0)) \\ &> \alpha. \end{aligned}$$

So,  $D(x_0)$  cannot satisfy (2). Sum up the above proof,  $A^*(x)$  is the supremum consisting of fuzzy sets in  $\mathcal{F}(X)$  to satisfy (2).

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